

# Technical Appendix to COMPETITIVE BIDDING IN AUCTIONS WITH PRIVATE AND COMMON VALUES\*

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## Appendix: Proofs

Below we shall need the following properties implied by logconcavity.

LEMMA A1. *If  $f(x, y)$  is logconcave then so are its marginals  $f(x)$  and  $f(y)$ .*

*Proof.* See An (1998).

COROLLARY A1. *The density  $f_s(s)$  is logconcave in  $s$ .*

*Proof.* The joint density of  $v$  and  $s$  is given by  $f(v, s) = f_v(v) f_s(v/n - s)$ . It is straightforward to verify that under Assumption 2 the Hessian of  $\log [f(v, s)]$  is negative semi-definite, i.e.  $f(v, s)$  is logconcave. Lemma A1 thus implies that  $f_s(s)$  is logconcave. Q.E.D.

LEMMA A2. *The hazard rate,  $h(x) = f(x)/[1-F(x)]$ , of a logconcave density,  $f(x)$ , is everywhere non-decreasing.*

*Proof.* Let  $f(x)$  be logconcave then

$$f'(x)[1 - F(x)] = \int_x^{x_H} f'(x)f(t)dt \geq \int_x^{x_H} f'(t)f(x)dt \geq -f(x)^2,$$

where the first inequality follows since logconcavity of  $f$  implies that  $f'(x)/f(x)$  is non-increasing in  $x$ , so  $f'(x)/f(x) \geq f'(t)/f(t)$  for all  $t \geq x$ . It is straightforward to rewrite the above equation as:  $\{f(x)/[1-F(x)]\}' \geq 0$ . Q.E.D.

LEMMA A3. *Let  $X$  and  $Y$  be random variables. The conditional expectation  $E(Y|X \leq x)$  satisfies*

$$\frac{\partial}{\partial x} E(Y|X \leq x) = [E(Y|X = x) - E(Y|X \leq x)] \frac{f_X(x)}{F_X(x)}.$$

*Proof.* Write the conditional expectation  $E(Y|X \leq x)$  as

$$E(Y|X \leq x) = \int_{x_L}^x E(Y|X = t) \frac{f_X(t)}{F_X(x)} dt.$$

Differentiating the right side with respect to  $x$  yields the desired result. Q.E.D.

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*Proof of Lemma 1.* Since  $f(v|s) = f_v(v)f_c(v/n-s)/f_s(s)$ , logconcavity of  $f_c(\cdot)$  is equivalent to

$$\partial_v \partial_s \log[f(v|s)] \geq 0.$$

This means that  $\partial_s \log[f(v|s)] = \partial_s f(v|s)/f(v|s)$  is increasing in  $v$ , or,  $\partial_s f(v_1|s)/f(v_1|s) \geq \partial_s f(v_2|s)/f(v_2|s)$  for all  $v_1 \geq v_2$ . This inequality can be rewritten as  $\partial_s[f(v_1|s)/f(v_2|s)] \geq 0$ . Hence, for all  $s_1 \geq s_2$  we have  $f(v_1|s_1)/f(v_2|s_1) \geq f(v_1|s_2)/f(v_2|s_2)$ .<sup>1</sup> Since this is true for all  $v_1 \geq v_2$ , it remains true if we integrate  $v_1$  from  $v$  to  $v_H$  and integrate  $v_2$  from  $v_L$  to  $v$ , which yields  $[1-F(v|s_1)]/F(v|s_1) \geq [1-F(v|s_2)]/F(v|s_2)$ , or,  $F(v|s_1) \leq F(v|s_2)$ . In other words,  $F(v|s_1)$  first-order stochastically dominates  $F(v|s_2)$ , which implies  $E(v|s_1) \geq E(v|s_2)$  for all  $s_1 \geq s_2$ . So  $E(v|s = x)$  is non-decreasing in  $x$ . Lemma A3 applied to  $X = s$ ,  $Y = v$  yields:  $\partial_x E(v|s \leq x) = [E(v|s = x) - E(v|s \leq x)]f_s(x)/F_s(x)$ , which is non-negative since  $E(v|s = x)$  is non-decreasing in  $x$ . The proofs for the conditional expectations of the cost are similar. Q.E.D.

Below we use the following notation. The conditional density of  $y_1$  given  $y_1 \leq s$  is:  $f_{y_1}(x|y_1 \leq s) = f_{y_1}(x)/F_{y_1}(s) = (n-1)f_s(x)F_s(x)^{n-2}/F_s(s)^{n-1}$  for  $x \leq s$ , the corresponding distribution function is  $F_{y_1}(x|y_1 \leq s) = F_s(x)^{n-1}/F_s(s)^{n-1}$ .

*Proof of Proposition 1.* An equivalent way to write (2) is:  $B(x) = (n-1)/n E(v|s \leq x) + E(y_1|y_1 \leq x)$  and its derivative is easily derived from Lemma A3:  $B'(x) = (n-1)/n [E(v|s = x) - E(v|s \leq x)] f_s(x)/F_s(x) + [x - E(y_1|y_1 \leq x)] f_{y_1}(x)/F_{y_1}(x)$ . The last term is positive (the event  $y_1 = x$  occurs with probability zero) and the first term is non-negative by Lemma 1. Hence,  $B(\cdot)$  is increasing. Suppose that bidders 2, ...,  $n$  bid according to (2). Bidder 1's expected profit is:

$$\pi_1^e(b) = \left\{ s_1 + \frac{n-1}{n} E[v|s \leq B^{-1}(b)] - b \right\} F_s^{n-1}[B^{-1}(b)].$$

From Lemma A3, the derivative of the expected profit with respect to  $b$  evaluated at  $b = B(x)$  is:

$$\pi_1^{e'}[B(x)] = \frac{f_{y_1}(x)}{B'(x)} \left[ s_1 + \frac{1}{n} E(v|s = x) + \frac{n-2}{n} E(v|s \leq x) - B(x) - B'(x) \frac{F_{y_1}(x)}{f_{y_1}(x)} \right].$$

Using the expression for  $B'$  derived above it is straightforward to verify that the term in the brackets equals  $s_1 - x$ . Together with monotonicity of  $B(\cdot)$  this shows  $B(s_1)$  is bidder 1's unique optimal bid. Q.E.D.

*Proof of Proposition 2.* Note that (3) is equivalent to:  $B(x) = x + (1/n) E(v|s = x) + [(n-2)/n] E(v|s \leq x)$ . The last two terms are non-decreasing in  $x$  (Lemma 1), so  $B(\cdot)$  is increasing. Assume that bidders 2, ...,  $n$  bid according to (3). Bidder 1's expected payoff is:

$$\pi_1^e(b) = \left\{ s_1 + \frac{n-1}{n} E[v|s \leq B^{-1}(b)] - \int_{s_L}^{B^{-1}(b)} B(x) f_{y_1}[x|s \leq B^{-1}(b)] dx \right\} F_s^{n-1}[B^{-1}(b)].$$

The second term in the brackets on the right side represents the sum of the values of bidders 2 to  $n$  given that their surpluses are less than  $B^{-1}(b)$ . An equivalent way of writing this term is by choosing one of them to have the highest surplus,  $y_1 \leq B^{-1}(b)$ , while the  $n-2$  other have surpluses less than or equal to  $y_1$ . The expected profit then becomes

<sup>1</sup> In other words,  $f(v|s)$  satisfies the monotone likelihood property, see Milgrom (1981).

$$\begin{aligned}\pi_1^e(b) &= \int_{s_L}^{B^{-1}(b)} \left[ s_1 + \frac{1}{n} E(v|s=x) + \frac{n-2}{n} E(v|s \leq x) - B(x) \right] f_{y_1}(x) dx \\ &= \int_{s_L}^{B^{-1}(b)} (s_1 - x) dF_{y_1}(x).\end{aligned}$$

Together with monotonicity of  $B^{-1}(\cdot)$  this proves that the unique optimal bid is  $B(s_1)$ . Q.E.D.

*Proof of Proposition 3.* Note that each  $B_k$  is strictly increasing in  $x$ . Suppose bidders  $2, \dots, n$  bid according to (4). When bidder 1 wins the auction her expected profit is:  $s_1 + 1/n \sum_{j=1}^{n-1} E(v|s = s_{j+1}) - B_{n-2}(s_2)$ , where the  $s_j$  are the realisations of the others' surpluses arranged in increasing order, i.e.  $s_2 \geq \dots \geq s_n$ . Using the definition of  $B_{n-2}$  the expected payoff can be written as:  $s_1 + (1/n) E(v|s = s_2) - (2/n) E(v|s = s_2) + E(d|s = s_2) = s_1 - s_2$ . So bidder 1's expected profit is positive only when she has the highest surplus, and using  $B(\cdot)$  she wins iff  $s_1 = Y_1$ . Hence,  $B(\cdot)$  is the optimal bidding strategy for player 1. Q.E.D.

*Proof of Proposition 4.* Expected total surplus equals the winner's valuation of the commodity:

$$W = \int_{s_L}^{s_H} \left[ x + \frac{n-1}{n} E(v|s \leq x) \right] f_{Y_1}(x) dx,$$

so  $W = E(s|s = Y_1) + [(n-1)/n] E(v|s \leq Y_1) = (1/n) E(v|s = Y_1) + [(n-1)/n] E(v|s \leq Y_1) - E(d|s = Y_1)$ . The first two terms represent the expected value of the commodity given that one bidder has the highest surplus and the others have lower surpluses: this is just  $E(V)$ . Hence  $W = E(V) - E(d|s = Y_1)$ .

The Envelope Theorem implies that the derivative of the equilibrium profit  $\pi^*(s) = \pi^e[B(s)]$  with respect to a bidder's surplus  $s$  equals the equilibrium probability of winning. When surpluses are i.i.d. across bidders the equilibrium probability of winning is simply  $F_s(s)^{n-1}$  and the equilibrium profit is therefore  $\pi^*(s) = \int_{s_L}^s F_s(x)^{n-1} dx$ . This is the expected equilibrium profit of a bidder with surplus  $s$  in any of the three auctions discussed above. The *ex ante* expected equilibrium *payoffs* accrue to the winning bidder, so

$$\pi_{\text{winner}} = n \int_{s_L}^{s_H} \int_{s_L}^s F_s^{n-1}(x) dx dF_s(s) = n \int_{s_L}^{s_H} [1 - F_s(x)] F_s^{n-1}(x) dx,$$

where the last expression is obtained by changing the order of integration. Partially integrating this expression yields:  $\pi_{\text{winner}} = E(Y_1) - E(Y_2)$ . Finally, revenue is the difference between total surplus and the winner's profit:  $R = W - \pi_{\text{winner}} = W - [E(Y_1) - E(Y_2)]$ . Q.E.D.

*Proof of Proposition 5.* Let  $\alpha_1 > \alpha_2 > 0$ . Suppose, without loss of generality, that  $s_1(\alpha_1) = \max_{j=1, \dots, n} s_j(\alpha_1)$  and  $s_k(\alpha_2) = \max_{j=1, \dots, n} s_j(\alpha_2)$ . There are two cases to be considered. If  $k = 1$ , the costs are the same for both  $\alpha$ s. If  $k \neq 1$ , we have:  $\alpha_1 v_1/n - c_1 > \alpha_1 v_k/n - c_k$  and  $\alpha_2 v_k/n - c_k > \alpha_2 v_1/n - c_1$  (ties occur with probability zero). Dividing the first inequality by  $\alpha_1$  and the second by  $\alpha_2$  and adding the result yields:  $c_1/\alpha_1 + c_k/\alpha_2 < c_1/\alpha_2 + c_k/\alpha_1$ , or, equivalently,  $c_k(1/\alpha_2 - 1/\alpha_1) < c_1(1/\alpha_2 - 1/\alpha_1)$ . Hence  $c_1 > c_k$  when  $k \neq 1$ . In other words, for some realisations of the value and cost signals, the winner's cost is less for  $\alpha = \alpha_2$  than for  $\alpha = \alpha_1$ , while it is the same for other realisations. *A fortiori*, the expected winner's

cost (i.e. the winner's cost averaged over all possible realisations of the value and cost signals) is less for  $\alpha = \alpha_2$ . So an increase in  $\alpha$  lowers total expected surplus. Next we show that an increase in  $\alpha$  does not lower bidders' expected profits. Note that the distribution function,  $F_{s(\alpha)}(x)$ , of  $s(\alpha) = \alpha v/n - c$  is given by

$$F_{s(\alpha)}(x) = \int_{v_L}^{v_H} f_v(v)[1 - F_c(\alpha v/n - x)]dv,$$

from which it follows that  $dF_{s(\alpha)}(x)/d\alpha = -1/n E[v|s(\alpha) = x]f_{s(\alpha)}(x) \leq 0$ , i.e. the distribution  $F_{s(\alpha)}(\cdot)$  is stochastically increasing in  $\alpha$ . The difference between the first and second order statistic is:

$$E[Y_1(\alpha)] - E[Y_2(\alpha)] = n \int_{s_L(\alpha)}^{s_H(\alpha)} F_{s(\alpha)}^{n-1}(x)[1 - F_{s(\alpha)}(x)]dx.$$

Differentiating this expression with respect to  $\alpha$  yields

$$\frac{d\{E[Y_1(\alpha)] - E[Y_2(\alpha)]\}}{d\alpha} = \frac{1}{n} \{E[v|s(\alpha) = Y_1(\alpha)] - E[v|s(\alpha) = Y_2(\alpha)]\}.$$

Lemma 1 together with  $Y_1(\alpha) > Y_2(\alpha)$  ensures that the right side is non-negative. Since bidders are no worse off and total expected surplus is lower, the seller's revenue has to be less. Q.E.D.

*Proof of Lemma 2.*

$$E_n(c|s = Y_1) = \int_{s_L}^{s_H} E(c|s = x)dF_s^n(x),$$

and

$$E_{n+1}[c|s = Y_1'(\alpha_0)] = \int_{s_L}^{s_H} E(c|s = x)dF_s^{n+1}(x).$$

Since  $F^{n+1}(\cdot)$  first-order dominates  $F^n(\cdot)$  and  $E(c|s = x)$  is non-increasing in  $x$ , we have  $E_{n+1}[c|s = Y_1'(\alpha_0)] \leq E_n(c|s = Y_1)$ . Q.E.D.